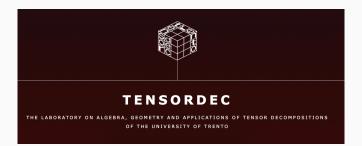
TensorDay 2023 Inaugural day of the TensorDec Laboratory

21 November 2023





Since 2021, the TensorDec Laboratory is a recipient for activities of research, teaching, mentoring and networking about algebra, geometry and applications of tensor decompositions



tensordec.maths.unitn.it





TensorDec Laboratory: who we are



Valentina Amitrano Dario Antolini



TensorDec Laboratory: who we are



Why should you knock at those doors? What are tensors and what do we do with them?



Tensors are multidimensional boxes for organizing numbers, much like matrices are boxes for organizing numbers in two ways.



Like matrices, tensors are extremely useful and versatile.

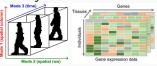
They can be viewed as:

- \bullet Tables \rightarrow Boxes of numbers with a certain order.
- Linear \rightarrow Multilinear maps.
- Elements of the Matrix \rightarrow Tensor space.



Tensors: as boxes of numbers

Multidimensional numerical arrays:



Used to store tons of data

extract meaningful information



Algorithms for Tensor Decompositions (come to our courses/masterclasses):

- Tucker Decomposition: $\mathcal{U} = (U_1, U_2, U_3)\mathcal{C}$:
- Rank Decomposition (parafac/candecomp/cpd):

$$\mathcal{U} = \sum_{i=1}^{r} u_{i,1} \otimes u_{i,2} \otimes u_{i,3}$$
:





A matrix $M \in W \otimes V$ is equivalent to a linear map $f_M : V \to W$. Multilinear maps

$$\begin{aligned} \mathsf{MaMu}_{2} : & (\mathbb{C}^{2,2})^{\times 2} & \to & \mathbb{C}^{2,2} \\ & \left(\begin{pmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{pmatrix}, \begin{pmatrix} b_{11} & b_{12} \\ b_{21} & b_{22} \end{pmatrix} \right) & \mapsto & \begin{pmatrix} a_{11}b_{11} + a_{12}b_{21} & a_{11}b_{12} + a_{12}b_{22} \\ a_{21}b_{11} + a_{22}b_{21} & a_{21}b_{12} + a_{22}b_{22} \end{pmatrix} \\ & \mathbb{C}^{2,2} \end{pmatrix}^{\otimes 3} \ni \mathsf{MaMu}_{2} = (\mathsf{a}_{11} \otimes \mathsf{b}_{11} + \mathsf{a}_{12} \otimes \mathsf{b}_{21}) \otimes \mathsf{c}_{11} + (\mathsf{a}_{11} \otimes \mathsf{b}_{12} + \mathsf{a}_{12} \otimes \mathsf{b}_{22}) \otimes \mathsf{c}_{12} \\ & \quad + (\mathsf{a}_{21} \otimes \mathsf{b}_{11} + \mathsf{a}_{22} \otimes \mathsf{b}_{21}) \otimes \mathsf{c}_{21} + (\mathsf{a}_{21} \otimes \mathsf{b}_{12} + \mathsf{a}_{22} \otimes \mathsf{b}_{22}) \otimes \mathsf{c}_{12} \end{aligned}$$

A decomposition of MaMu₂ highlights an algorithm.

The rank decomposition highlights the best one. Rank = Complexity \rightarrow Complexity Theory.

If the algorithm is encoded in a Quantum Circuit: Complexity = Entanglement of the state \rightarrow Quantum Information.



Points of a tensor space: matrix $M \in W \otimes V$, tensor $U \in V_1 \otimes \cdots \otimes V_d$



This perspective allows to study tensors by their common properties, using global techniques like algebraic-geometric methods (both symbolic and numerical) to extract shared information among them.

- Mima Stanojkovski: "Tensors in finite group theory".
- Alessandro Oneto: "Tensors into algebraic statistical models".
- Elisa Postinghel: "Tensors and polynomial interpolation".
- Alessandra Bernardi: "Symbolic algorithms for Tensor Decomposition".
- Edoardo Ballico: "Algebraic Geometry aspects of Tensors".



We mentor theses

Algebra





Mima Stanojkovski Alessandra Bernardi

- Commutative and non commutative algebra, Group theory, Scheme theory.
- Interdisciplinary and Industrial: Quantum physics, Data science.

Algebraic Geometry



Alessandro Oneto Elisa Postinghel

Edoardo Ballico

- Birational geometry, Computational geometry, Teaching.
- Interdisciplinary and Industrial: Data science, Algebraic statistics.



Courses (Mathematics for Data Science, Q@TN):

- Tensor Decompositions for Big Data Analysis (A. Bernardi)
- Geometry and Topology for Data Analysis (A. Oneto)

We organize annual Masterclasses taught by international guests:





TensorDec Laboratory: working with us - Networking

PhD opportunities (more than the standard call):

- Industrial PhD (just closed)
- Transdisciplinary Doctoral Program Q@TN
- Horizon: Marie Curie Double Degree (IT & abroad):

TENORS Tensor modEliNg, geOmetRy and optimiSation Marie Skłodowska-Curie Doctoral Network



Tensors are nowadays ubiquitous in many domains of applied mathematics, computer science, signal processing, data processing, machine learning and in the emerging area of quantum computing. TENOR5 aims at fostering cutting-edge research in tensor sciences, stimulating interdisciplinary and intersectoriality knowledge developments between algebraists, geometers, computer scientists, numerical analysts, data analysts, physicits, quantum scientists, and industrial actors facing real-life tensor-based problems.

Partners:

- Inria, Sophia Antipolis, France (B. Mourrain, A. Mantzaflaris)
- 2 CNRS, LAAS, Toulouse, France (D. Henrion, V. Magron, M. Skomra)
- 8 NWO-I/CWI, Amsterdam, the Netherlands (M. Laurent)
- Univ. Konstanz, Germany (M. Schweighofer, S. Kuhlmann, M. Michałek)
- 6 MPI, Leipzig, Germany (B. Sturmfels, S. Telen)
- 6 Univ. Tromsoe, Norway (C. Riener, C. Bordin, H. Munthe-Kaas)
- O Univ. degli Studi di Firenze, Italy (G. Ottaviani)
- Univ. degli Studi di Trento, Italy (A. Bernardi, A. Oneto, I. Carusotto)
- OTU, Prague, Czech Republic (J. Marecek)
- 10 ICFO, Barcelona, Spain (A. Acin)
- Artelys SA, Paris, France (M. Gabay)

Associate partners:

- Quandela, France
- 2 Cambridge Quantum Computing, UK.
- Bluetensor, Italy.
- Arva AS, Norway.
- 6 HSBC Lab., London, UK.

15 PhD positions (2024-2027)

(recruitment expected around Oct. 2024)

Scientific coord: B. Mourrain Adm. manager: Linh Nguyen



TensorDec Laboratory: working with us - Networking

Post Doc opportunities: We are members of the Italian Network for Applied and Birational Algebraic Geometry (INABAG)





https://sites.google.com/unitn.it/inabag/



Let $R_d := \mathbb{R}[x_1, x_2]_d$ be the space of polynomials of degree $\leq d$. Let $S = \{p_1, \dots, p_s\}$ be a set of distinct points in \mathbb{R}^2 .



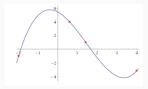
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Polynomial interpolation problems:

• Simple points:

$$\mathcal{L}_d(S) = \{f \in R : f(p_i) = 0, \forall i = 1, \dots, s\} \subseteq R$$







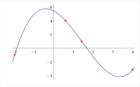
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• Double points:

 $\mathcal{L}_d(2S) = \{F \in R_d : f(p_i) = \frac{\partial}{\partial x_1} f(p_j) = \frac{\partial}{\partial x_2} f(p_j) = 0, \forall j = 1, \dots, s\} \subseteq R$



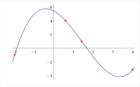
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• Multiple points:

 $\mathcal{L}_d(mS) = \dots$



$$R_d := \mathbb{R}[x_1, \ldots, x_n]_d$$

$$S = \{p_1, \ldots, p_s\} \subset \mathbb{R}^n$$

 \downarrow passing to algebraically closed fields

 $R_{\mathbb{C},d} := \mathbb{C}[x_1, \ldots, x_n]_d$ $S = \{p_1, \ldots, p_s\} \subset \mathbb{C}^n.$



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 \downarrow homogenising/projectivising

 $R^{\text{homog}}_{\mathbb{C},d} := \mathbb{C}[x_0, x_1 \dots, x_n]_d$ $S = \{p_1, \dots, p_s\} \subset \mathbb{P}^n_{\mathbb{C}}.$



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• Double point polynomial interpolation:

$$\mathcal{L}_d(2S) = \{F \in R_{\mathbb{C},d}^{\text{homog}} : \frac{\partial}{\partial x_i} F(p_j) = 0, \forall i = 1, \dots, s, \forall j = 0 \dots, n\}$$



Slogan

Double point polynomial interpolation problems

 $\stackrel{\hbox{Dimensionality of secant varieties to}}{\longleftrightarrow} \ \ \, \mbox{varieties of rank one symmetric tensors}$

 $\sigma_t(X_d)$



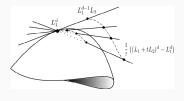
Slogan

Double point polynomial interpolation problems

Dimensionality of secant varieties to varieties of rank one symmetric tensors

 $\sigma_t(X_d)$

- Symmetric tensors: e.g. $\operatorname{Sym}^2(\mathbb{C}^n) = \{x \otimes y - y \otimes x : x, y \in \mathbb{C}^n\} \subset \mathbb{C}^n \otimes \mathbb{C}^n$ $\operatorname{Sym}^d(\mathbb{C}^n) = R^{homog}_{\mathbb{C},d}$ i.e. degree-*d* polynomials
- $X_d \subseteq \mathbb{P}(\operatorname{Sym}^d(\mathbb{C}^n))$ the Variety of rank-1 symmetric tensors
- $\sigma_t(X_d)$ is the *t*-secant variety to X_d







- **Computations** (isomorphism testing, automorphisms, ...)
- Enumeration (number of groups with given properties, number of groups with a shared quotient, ...)
- Global techniques for local objects (groups as points of a scheme).



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Question. Is it restrictive? If so, how restrictive is it?



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Finite groups are understood via their simple composition factors and via their Sylow *p*-subgroups. The first being classified, we look into the second.



Let p be a prime and $G_p = \langle g_1, g_2, g_3, h_1, h_2, h_3, z_1, z_2, z_3 \mid \text{relations} \rangle$ where:



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- *z*₁, *z*₂, *z*₃ central,
- $\langle g_1, g_2, g_3
 angle$ and $\langle h_1, h_2, h_3
 angle$ abelian,
- $[g_1, h_1] = z_1$, $[g_1, h_2] = z_2$, $[g_1, h_3] = 1$
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$$M = \begin{pmatrix} y_1 & y_2 & 0 \\ y_2 & -y_3 & y_1 \\ 0 & y_1 & -y_3 \end{pmatrix}$$



$$\det(M) = y_1 y_3^2 - y_1^3 + y_3 y_2^2$$



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... and G_p is a *p*-realization of an object defined over \mathbb{Z} ! ("Globalization")







Question. What if I had another curve in the plane?

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Question. What are the isomorphisms of G_p ? (Enumeration)

 $|\operatorname{Aut}(G_p)| = (\text{polynomial in } p) \times (\text{non-quasipolynomial}).$



Let X_1, \ldots, X_d discrete random with states $X_i \in [n_i] = \{1, \ldots, n_i\}$. Then we consider the tensor of **joint probabilities**: $T_{x_1, \ldots, x_d} = P(X_1 = x_1, \ldots, X_d = x_d)$.



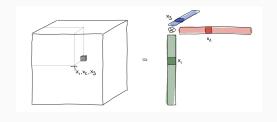
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If we assume the variables to be **independent** then

$$P(X_1 = x_1, \ldots, X_d = x_d) = P(X_1 = x_1) \cdots P(X_d = x_d),$$

namely, if we let $v_i = (P(X_i = 1), ..., P(X_i = n_i))$,

 $T = v_1 \otimes v_2 \otimes \cdots \otimes v_d$





Tensors in Statistical Models

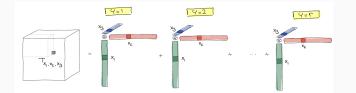
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If we assume the variables to be independent conditionally to $Y \in [r]$ then

$$P(X_1 = x_1, \dots, X_d = x_d) = \sum_{y=1}^r P(Y = y) P(X_1 = x_1, \dots, X_d = x_d | Y = y)$$
$$= \sum_{y=1}^r P(Y = y) P(X_1 = x_1 | Y = y) \cdots P(X_d = x_d | Y = y).$$

namely, if we let $v_i^{(y)} = (P(X_i = 1 | Y = y), \dots, P(X_i = n_i | Y = y))$, then

$$T = \sum_{y=1}^{r} \lambda_y v_1^{(y)} \otimes v_2^{(y)} \otimes \cdots \otimes v_d^{(y)}$$





statistical model which depends polynomially on its parameters; namely

a polynomial map $\varphi: \mathcal{P} \longrightarrow \mathcal{M} \subset \Delta_N \subset \mathbb{R}^N$

 $\mathcal{P} = \text{parameter space} \quad \mathcal{M} = \varphi(\mathcal{P}) \subset \Delta_N = \left\{ (p_0, \dots, p_N) \in \mathbb{R}^N \ : \ \frac{p_0 + \dots + p_N = 1}{p_i \geq 0} \right\}$



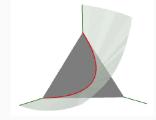
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Example: Two throws of a biased coin C. Let $Z \in [2]$ count the number of heads.

If P(C = H) = a and $p_i = P(Z = i)$, then $a \mapsto (p_0, p_1, p_2) = ((1 - a)^2, 2a(1 - a), a^2).$ Then, $\mathcal{M} = \{(p_0, p_1, p_2) \in \Delta_2 : p_1^2 - 4p_0p_2 = 0\}.$





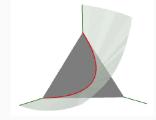
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Questions.

- What is the *dimension* of \mathcal{M} ?
- What are the defining *equations* and *inequalities* of \mathcal{M} ?
- Given y ∈ M, how φ⁻¹(y) look like?
 Is the model *identifiable*, i.e., the general fiber φ⁻¹(y) is a singleton?



Gaussian Models

Given a density function $f : \mathbb{R}^n \to \mathbb{R}$ for a random vector $X = (X_1, \ldots, X_n)$, its **moments** are

$$m_{i_1,\ldots,i_n} = \int_{\mathbb{R}^n} x_1^{i_1}\cdots x_n^{i_n} f(x_1,\ldots,x_n) dx_1\cdots dx_n$$

Example. For n = 1, if $X \sim \mathcal{N}(\mu, \sigma^2)$, then

$$m_1 = \mu, \ m_2 = \mu^2 + \sigma^2, \ m_3 = \mu^3 + 3\mu\sigma^2, \ m_4 = \mu^4 + 6\mu^2\sigma^2 + 3\sigma^4, \dots$$

In general, moments of Gaussian models are polynomials in the parameters!



Gaussian Models

Given $X_i \sim \mathcal{N}(\mu_i, \Sigma)$, i = 1, ..., m, and their *mixture* $Y = \lambda_1 X_1 + ... + \lambda_m X_m$, with $\lambda_1 + ... + \lambda_m = 1$, then

$$\varphi_d: (\lambda_1,\ldots,\lambda_m,\ldots,\mu_i,\Sigma_i,\ldots)\mapsto (m_{d0\cdots 0},m_{d-1,1,0\cdots 0},\ldots,m_{00\cdots d}).$$

This defines the degree-d moment variety of mixture of Gaussian models.

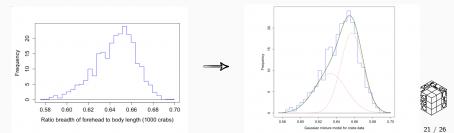


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Gaussian Models

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$$\varphi_d: (\lambda_1,\ldots,\lambda_m,\ldots,\mu_i,\Sigma_i,\ldots)\mapsto (m_{d0\cdots 0},m_{d-1,1,0\cdots 0},\ldots,m_{00\cdots d}).$$

This defines the degree-*d* moment variety of mixture of Gaussian models. Method of moments (Pearson's crabs, 1903)



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This defines the degree-d moment variety of mixture of Gaussian models.

Note:

 $\operatorname{im}(\varphi_d)$ can be looked inside the space of degree-d multivariate polynomials

$$\bar{m}_d = \sum_{i_1,\ldots,i_n} m_{i_1,\ldots,i_n} x_1^{i_1} \cdots x_n^{i_n}$$

If $\ell_i = \mu_i \cdot (x_1, ..., x_n)^T$ and $q_i = (x_1, ..., x_n) \Sigma_i (x_1, ..., x_n)^T$, then $\bar{m}_1 = \lambda_1 \ell_1 + ... + \lambda_m \ell_m$, $\bar{m}_2 = \lambda_1 (\ell_1^2 + q_1) + ... + \lambda_m (\ell_m^2 + q_m)$ $\bar{m}_3 = \lambda_1 (\ell_1^3 + 3\ell_1 q_1) + ... + \lambda_m (\ell_m^3 + 3\ell_m q_m)$, ...



Tensors under group actions

Main aim: study geometric properties of secant varieties, e.g.

- identifiability of points = uniqueness in recovering data ;
- singularity of points = unfeasibility in computations .



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Toy case: $\bigwedge^2 \mathbb{C}^n := \{A \in Mat_{n \times n}(\mathbb{C}) \text{ skew-symm.}\}$

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$$\operatorname{Gr}_{2,n} := \left\{ [A] \in \mathbb{P} \left(\bigwedge^2 \mathbb{C}^n \right) \mid \operatorname{rk}(A) = 2 \right\};$$

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The group $GL_n(\mathbb{C})$ acts on $\bigwedge^2 \mathbb{C}^n$ and leaves each $\sigma_r(Gr_{2,n})$ invariant. $GL_n(\mathbb{C})$ also preserves identifiability and singularity



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 \implies Enough to check them only on representatives of $GL_n(\mathbb{C})$ -orbits.



Let G be a "nice" group. $= \operatorname{GL}_n(\mathbb{C})$ Let V^G be a vector space on which G acts without invariant proper vector subspaces. $= \bigwedge^d \mathbb{C}^n$ space of skew-symmetric tensors in $\bigwedge^d \mathbb{C}^n$ Let $v_0 \in V^G$ be a "special" vector. $= e_1 \land \ldots \land e_d$



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- The orbit X := G · [v₀] is the variety of rank−1 tensors in P(V^G);
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Alice: Well, they are cool, that's why!

Bob: For instance, you find them in Quantum Physics ($Gr_{d,n}$ as "simple" fermions), in Quantum Chemistry ($Gr_{d,n}$ as quantum states of *d* electrons in *n* orbitals), in Quantum Information ("isotropic" Grassmannians parametrize abelian groups of observables in the Clifford group), ...



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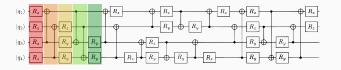
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Goal: Find the best quantum gate decomposition $U = U_L \dots U_2 \cdot U_1$





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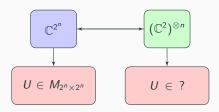


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Entangling gates are quasi-multilinear maps

$$|\psi'
angle = \sum_{k=1}^{s} (M_k^1, \dots, M_k^n) \cdot |\psi
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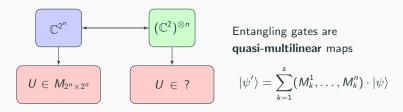


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Goal: Find the best quantum gate decomposition can be viewed as find the rank decomposition of U

$$U=\sum_{k=1}^{s}M_{k}^{1}\otimes\cdots\otimes M_{k}^{n}$$



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